# Computation of Moral-Hazard Problems* 

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#### Abstract

This paper studies computational aspects of moral-hazard problems. In particular, we consider deterministic contracts as well as contracts with action and/or compensation lotteries, and formulate each case as a mathematical program with equilibrium constraints (MPEC). We investigate and compare solution properties of the MPEC approach to that of the linear programming (LP) approach with lotteries. We propose a hybrid procedure that combines the best features of the both. The hybrid procedure obtains a solution that is, if not global, at least as good as an LP solution. It also preserves the fast local convergence property by applying the SQP algorithm to MPECs. The numerical results on an example show that the hybrid procedure outperforms the LP approach in both computational time and solution quality in term of the optimal objective value.


## 1 Introduction

This paper studies mathematical programming approaches to solve moral-hazard problems. More specifically, we formulate moral-hazard problems with finitely many action choices, including the basic deterministic models and models with lotteries, as mathematical programs with equilibrium constraints (MPECs). One advantage of using an MPEC formulation is the size of resulting program, which is often a thousand times smaller than the linear programs derived from the LP lotteries approach [18, 19]. This feature makes the MPEC approach an

[^0]appealing alternative when solving a large-scale linear program is computationally infeasible because of limitations on computer memory or computing time.

The moral-hazard model studies the relationship between a principal (leader) and an agent (follower) in situations in which the principal can neither observe nor verify an agent's action. The model is formulated as a bilevel program, in which the principal's upper-level decision takes the agent's best response to the principal's decision into account. Bilevel programs are generally difficult mathematical problems and much research in the economics literature has been devoted to analyzing and characterizing solutions of the moral-hazard model (see Grossman and Hart [7] and the references therein). When the agent's set of actions is a continuum, an intuitive approach to simplify the model is to assume the agent's optimal action lies in the interior of the action set. One then can treat the agent's problem as an unconstrained maximization problem and replace it by the first-order optimality conditions. This is called the first-order approach in the economics literature. However, Mirrlees [12, 13] showed that the first-order approach may be invalid because the lower-level agent's problem is not necessarily a concave maximization program and that the optimal solution may fail to be unique and interior. Consequently, a sequence of papers [20, 8, 9] has developed conditions under which the first-order approach is valid. Unfortunately, these conditions are often more restrictive than is desirable.

In general, if the lower-level problem in a bilevel program is a convex minimization (or concave maximization) problem, one can then replace the lower-level problem by the firstorder optimality conditions, which are both necessary and sufficient, and reformulate the original bilevel problem as an MPEC. This idea is similar to the first-order approach to the moral-hazard problem with one notable difference: MPEC formulations include complementarity constraints. The first-order approach assumes that the solution to the agent's problem lies in the interior of the action set, and hence, one can treat the agent's problem as an unconstrained maximization problem. This assumption may also avoid issues associated with the failure of the constraint qualification at a solution. General bilevel programs do not assume an interior solution assumption. As a result, the complementarity conditions associated with the Karush-Kuhn-Tucker multipliers for inequality constraints would appear in the first-order optimality conditions for the lower-level program. MPECs also arise in many applications in engineering (e.g., transportation, contact problems, mechanical structure design) and economics (Stackelberg games, optimal taxation problems). One well-known theoretical difficulty with MPECs is that the standard constraint qualification such as linear independence constraint qualification and Mangasarian-Fromovitz constraint qualification fails at every feasible point. An extensive literature has developed to refine constraint qualifications and stationarity conditions for MPECs; see Scheel and Scholtes [22] and the references therein. We also refer to the two-volume monograph by Facchinei and Pang [2] for theory and applications of equilibrium problems and to the monographs by Luo et al. [11] and Outrata et al. [16] for more details on MPEC theory and applications.

The failure of the constraint qualification conditions means that the set of Lagrange multipliers is unbounded and that conventional numerical optimization software may fail to
converge to a solution. Economists have avoided these numerical problems by reformulating the moral-hazard problem as a linear program involving lotteries over a finite set of outcomes. See Townsend [24, 25] and Prescott [18, 19]. While this approach avoids the constraint qualification problems, it does so by restricting aspects of the contract, such as consumption, to a finite set of possible choices even though a continuous choice formulation would be economically more natural.

The purpose of this paper is twofold: (1) To introduce to the economics community the MPEC approach, or more generally, advanced equilibrium programming approaches, to the moral hazard problem; (2) To present an interesting and important class of incentive problems in economics to the mathematical programming community. Many incentive problems, such as contract design, optimal taxation and regulation, and multiproduct pricing, can be naturally formulated as an MPEC or an equilibrium problem with equilibrium constraints (EPEC) [23]. This greatly extends the applications of equilibrium programming to one of the most active research topics in economics in past three decades. The need for a global solution for these economical problems provides a motivation for optimization community to develop efficient global optimization algorithms for MPECs and EPECs.

The remainder of this paper is organized as follows. In the next section, we describe the basic moral-hazard model and formulate it as a mixed-integer nonlinear program and as an MPEC. In Section 3, we consider moral-hazard problems with action lotteries, with compensation lotteries, and with the combination of the both. We derive MPEC formulations for each of these cases. We also compare the properties of the MPEC approach and the LP lottery approach. In Section 5, we develop a hybrid approach that preserves the robustness of global solution from the LP approach and the fast local convergence of the MPEC approach. The numerical efficiency of the hybrid approach in both computational speed and robustness of the solution is illustrated in an example.

## 2 The Basic Moral-Hazard Model

### 2.1 The Deterministic Contract

We consider a moral-hazard model in which the agent chooses an action from a finite set $\mathcal{A}=\left\{a_{1}, \ldots, a_{M}\right\}$. The outcome can be one of $N$ alternatives. Let $\mathcal{Q}=\left\{q_{1}, \ldots, q_{N}\right\}$ denote the outcome space where the outcomes are dollar returns to the principal ordered from smallest to largest.

The principal can only observe the outcome, not the agent's action. However, the stochastic relationship between actions and outcomes, which is often called a production technology is common knowledge to both the principal and the agent. Usually, the production technology is exogenously described by the probability distribution function, $p(q \mid a)$, which presents the probability of outcome $q \in \mathcal{Q}$ occuring given that action $a$ is taken. We
assume $p(q \mid a)>0$ for all $q \in \mathcal{Q}$ and $a \in \mathcal{A}$; this is called the full-support assumption.
Since the agent's action is not observable to the principal, the payment to the agent is only based on the outcome observed by the principal. Let $\mathcal{C} \subset R$ be the set of all possible compensations.

Definition 1 A compensation schedule $c=\left(c\left(q_{1}\right), \ldots, c\left(q_{N}\right)\right) \in R^{N}$ is an agreement between the principal and the agent such that $c(q) \in \mathcal{C}$ is the payoff to the agent from the principal if outcome $q \in \mathcal{Q}$ is observed.

The agent's utility $u(x, a)$ is a function of the payment $x \in R$ received from the principal and of his action $a \in \mathcal{A}$. The principal's utility $w(q-x)$ is a function over net income $q-x$ for $q \in \mathcal{Q}$. We let $W(c, a)$ and $U(c, a)$ denote the expected utility to the principal and agent, respectively, of a compensation schedule $c \in R^{N}$ when the agent chooses action $a \in \mathcal{A}$, i.e.,

$$
\begin{align*}
W(c, a) & =\sum_{q \in \mathcal{Q}} p(q \mid a) w(q-c(q)),  \tag{1}\\
U(c, a) & =\sum_{q \in \mathcal{Q}} p(q \mid a) u(c(q), a) .
\end{align*}
$$

Definition 2 A deterministic contract (proposed by the principal) consists of a recommended action $a \in \mathcal{A}$ to the agent and a compensation schedule $c \in R^{N}$.

The contract has to satisfy two conditions to be accepted by the agent. The first condition is the participation constraint. It states that the contract must give the agent an expected utility no less than a required utility level $U^{*}$ :

$$
\begin{equation*}
U(c, a) \geq U^{*} \tag{2}
\end{equation*}
$$

The value $U^{*}$ represents the highest utility the agent can receive from other activities if he does not sign the contract.

Second, the contract must be incentive compatible to the agent; it has to provide incentives for the agent not to deviate from the recommended action. In particular, given the compensation schedule $c$, the recommended action $a$ must be optimal from the agent's perspective and maximize the agent's expected utility function. The incentive compatibility constraint is given as follows:

$$
\begin{equation*}
a \in \operatorname{argmax}\{U(c, a): a \in \mathcal{A}\} . \tag{3}
\end{equation*}
$$

For a given $U^{*}$, a feasible contract satisfies the participation constraint (2) and the incentive compatibility constraint (3). The objective of the principal is to find an optimal deterministic contract, a feasible contract that maximize his expected utility. A mathemat-
ical program of finding an optimal deterministic contract $\left(c^{*}, a^{*}\right)$ is:

$$
\begin{array}{ll}
\operatorname{maximize}_{(c, a)} & W(c, a) \\
\text { subject to } & U(c, a) \geq U^{*},  \tag{4}\\
& a \in \operatorname{argmax}\{U(c, a): a \in \mathcal{A}\} .
\end{array}
$$

Since there are only finitely many actions in $\mathcal{A}$, the incentive compatibility constraint (3) can be presented as the following set of inequalities:

$$
\begin{equation*}
U(c, a) \geq U\left(c, a_{i}\right), \quad \text { for } i=1, \ldots, M \tag{5}
\end{equation*}
$$

These constraints ensure that the agent's expected utility obtained from choosing the recommendation action is no worse than that of choosing other actions. Replacing (3) by the set of inequalities (5), we have an equivalent formulation of the optimal contract problem:

$$
\begin{array}{ll}
\operatorname{maximize}_{(c, a)} & W(c, a) \\
\text { subject to } & U(c, a) \geq U^{*} \\
& U(c, a) \geq U\left(c, a_{1}\right)  \tag{6}\\
& \vdots \\
& U(c, a) \geq U\left(c, a_{M}\right) \\
& a \in \mathcal{A}=\left\{a_{1}, \ldots, a_{M}\right\} .
\end{array}
$$

### 2.2 A Mixed-Integer Nonlinear Programming Formulation

The optimal contract problem (6) can be formulated as a mixed-integer nonlinear program. Associated with each action $a_{i} \in \mathcal{A}$, we introduce a binary variable $y_{i} \in\{0,1\}$. Let $y=$ $\left(y_{1}, \ldots, y_{M}\right) \in R^{M}$ and let $e_{M}$ denote the vector of all ones in $R^{M}$. To ease the notation, we define

$$
\begin{align*}
U(c) & =\left(U\left(c, a_{1}\right), \ldots, U\left(c, a_{M}\right)\right) \in R^{M} \\
W(c) & =\left(W\left(c, a_{1}\right), \ldots, W\left(c, a_{M}\right)\right) \in R^{M} \tag{7}
\end{align*}
$$

The mixed-integer nonlinear programming formulation for the optimal contract problem (6) is

$$
\begin{array}{ll}
\operatorname{maximize}_{(c, y)} & y^{\mathrm{T}} W(c) \\
\text { subject to } & y^{\mathrm{T}} U(c) \geq U^{*}, \\
& y^{\mathrm{T}} U(c) \geq U\left(c, a_{1}\right), \\
& \vdots  \tag{8}\\
& y^{\mathrm{T}} U(c) \geq U\left(c, a_{M}\right) \\
& e_{M}^{\mathrm{T}} y=1, \\
& y_{i} \in\{0,1\} \quad \forall i=1, \ldots, M
\end{array}
$$

The above problem has $|\mathcal{Q}|$ nonlinear variables, $|\mathcal{A}|$ binary variables, one linear constraint and $(|\mathcal{A}|+1)$ nonlinear constraints. To solve a mixed-integer nonlinear program, one can use MINLP [3], BARON [21] or other solvers developed for this class of programs. For (8), since the agent will choose one and only one action, the number of possible combinations on the binary vector $y$ is only $M$. One then can solve (8) by solving $M$ nonlinear programs with $y_{i}=1$, and $y_{j}=0$ in the $i$-th nonlinear program, as Grossman and Hart suggested in [7] for the case where the principal is risk averse. They further point out that each nonlinear program is a convex program if the agent's utility function $u(x, a)$ can be written as $G(a)+K(a) V(x)$, where (1) $V$ is real-valued, strictly increasing, concave function defined on some open interval $\mathcal{I}=(\underline{I}, \bar{I}) \subset R$; (2) $\lim _{x \rightarrow \underline{I}} V(x)=-\infty$; (3) $G, K$ are real-valued functions defined on $\mathcal{A}$ and $K$ is strictly positive; (4) $u(x, a) \geq u(x, \hat{a}) \Rightarrow u(\hat{x}, a) \geq u(\hat{x}, \hat{a})$, for all $a, \hat{a} \in \mathcal{A}$, and $x, \hat{x} \in \mathcal{I}$. The above assumption implies that the agent's preferences over income lotteries are independent of his action.

### 2.3 An MPEC Formulation

In general, a mixed-integer nonlinear program is a difficult optimization problem. Below, by considering a mixed-strategy reformulation of the incentive compatibility constraints for the agent, we can reformulate the optimal contract problem (6) as a mathematical program with equilibrium constraint (MPEC); see [11].

For $i=1, \ldots, M$, let $\delta_{i}$ denote the probability that the agent will choose action $a_{i}$. Then, given the compensation schedule $c$, the agent chooses a mixed strategy profile $\delta^{*}=$ $\left(\delta_{1}^{*}, \ldots, \delta_{M}^{*}\right) \in R^{M}$ such that

$$
\begin{equation*}
\delta^{*} \in \operatorname{argmax}\left\{\sum_{k=1}^{M} \delta_{k} U\left(c, a_{k}\right): e_{M}^{\mathrm{T}} \delta=1, \delta \geq 0\right\} . \tag{9}
\end{equation*}
$$

Observe that the agent's mixed-strategy problem (9) is a linear program, and hence, its optimality conditions are necessary and sufficient.

The following lemma states the relationship between the optimal pure strategy $a_{i}$ and the optimal mixed strategy $\delta^{*}$.

Lemma 1 Given a compensation schedule $\bar{c} \in R^{N}$, the agent's action $a_{i} \in \mathcal{A}$ is optimal to the problem (3) iff there exists an optimal mixed strategy profile $\delta^{*}$ to the problem (9) such that

$$
\begin{gathered}
\delta_{i}^{*}>0 \\
\sum_{k=1}^{M} \delta_{k}^{*} U\left(\bar{c}, a_{k}\right)=U\left(\bar{c}, a_{i}\right) \\
e_{M}^{\mathrm{T}} \delta^{*}=1, \quad \delta^{*} \geq 0
\end{gathered}
$$

Proof If $a_{i}$ is an optimal action of (3), then let $\delta^{*}=e_{i}$, where $e_{i}$ is the $i$-th column in an identity matrix of order $M$. It is easy to verify that all the conditions for $\delta^{*}$ are satisfied. Conversely, if $a_{i}$ is not an optimal solution of (3), then there exists an optimal action $a_{j}$ such that $U\left(\bar{c}, a_{j}\right)>U\left(\bar{c}, a_{i}\right)$. Let $\tilde{\delta}=e_{j}$. Then $\tilde{\delta}^{\mathrm{T}} U(c)=U\left(\bar{c}, a_{j}\right)=U\left(\bar{c}, a_{i}\right)>\delta^{* \mathrm{~T}} U(c)$. We have a contradiction.

Let $u \perp v$ indicate orthogonality of vectors $u$ and $v$, i.e., $u^{\mathrm{T}} v=0$. An observation following from Lemma 1 is stated below.

Lemma 2 Given a compensation schedule $c \in R^{N}$, a mixed strategy profile $\delta$ is optimal to the linear program (9) iff

$$
\begin{gather*}
0 \leq \delta \perp\left(\delta^{\mathrm{T}} U(c)\right) e_{M}-U(c) \geq 0  \tag{10}\\
e_{M}^{\mathrm{T}} \delta=1
\end{gather*}
$$

Proof. This follows from the optimality conditions and the strong duality theorem for the LP (9).

Substituting the incentive compatibility constraint (5) by the system (10) and replacing $W(c, a)$ and $U(c, a)$ by $\delta^{\mathrm{T}} W(c)$ and $\delta^{\mathrm{T}} U(c)$, respectively, we derive an MPEC formulation of the principal's problem (6):

$$
\begin{array}{ll}
\operatorname{maximize}_{(c, \delta)} & \delta^{\mathrm{T}} W(c) \\
\text { subject to } & \delta^{\mathrm{T}} U(c) \geq U^{*} \\
& e_{M}^{\mathrm{T}} \delta=1  \tag{11}\\
& 0 \leq \delta \perp\left(\delta^{\mathrm{T}} U(c)\right) e_{M}-U(c) \geq 0
\end{array}
$$

To illustrate the failure of constraint qualification at any feasible point of an MPEC, we consider the feasible region $\mathcal{F}_{1}=\left\{(x, y) \in R^{2} \mid x \geq 0, y \geq 0, x y=0\right\}$. At the point $(\bar{x}, \bar{y})=(0,2)$, the first constraint $x \geq 0$ and the third constraint $x y=0$ are binding. Vectors of the gradient of the binding constraints at $(\bar{x}, \bar{y})$ are $(1,0)$ and $(2,0)$, which are dependent. It is easy to verify that the gradient vectors of the binding constraints are indeed dependent at other feasible points.


Fig. 1. The feasible region $\mathcal{F}_{1}=\{(x, y) \mid x \geq 0, y \geq 0, x y=0\}$.

The following lemma states the relationship between the optimal solutions for the principalagent problems (6) and the corresponding MPEC formulation (11).

Theorem 1 If $\left(c^{*}, \delta^{*}\right)$ is an optimal solution for the MPEC (11), then $\left(c^{*}, a_{i}^{*}\right)$, where $i \in$ $\left\{j: \delta_{j}^{*}>0\right\}$, is an optimal solution for the problem (6). Conversely, if ( $c^{*}, a_{i}^{*}$ ) is an optimal solution for the problem (6), then $\left(c^{*}, e_{i}\right)$ is an optimal solution for the MPEC (11).

Proof The statement follows directly from Lemma 2.
The MPEC (11) has $(|\mathcal{Q}|+|\mathcal{A}|)$ variables, 1 linear constraint, 1 nonlinear constraint, and $|\mathcal{A}|$ complementarity constraints. Hence, the size of the problem grows linearly in the number of the outcomes and actions. As we will see in Section 3.4, this feature is the main advantage of using the MPEC approach comparing to the LP lotteries approach.

## 3 Moral-Hazard Problems with Lotteries

In this section, we study moral-hazard problems with lotteries. In particular, we consider action lotteries, compensation lotteries, and the combination of the both. For each case, we first give definitions for the associated lotteries and then derive the nonlinear programming or MPEC formulation.

### 3.1 The Contract with Action Lotteries

Definition 3 A contract with action lotteries is a probability distribution over actions, $\pi(a)$, and a compensation schedule $c(a)=\left(c\left(q_{1}, a\right), \ldots, c\left(q_{N}, a\right)\right) \in R^{N}$ for all $a \in \mathcal{A}$. The compensation schedule $c(a)$ is an agreement between the principal and the agent such that $c(q, a) \in \mathcal{C}$ is the payoff to the agent from the principal if outcome $q \in \mathcal{Q}$ is observed and the action $a \in \mathcal{A}$ is recommended by the principal.

In the definition of a contract with action lotteries, the compensation schedule $c(a)$ is contingent on both the outcome and the agent's action. Given this definition, one might raise the following question: if the principal can only observe the outcome, not the agent's action, is it reasonable to have the compensation schedule $c(a)$ contingent on the action chosen by the agent? After all, the principal does not know what action is implemented by the agent. One economic justification is as follows. Suppose that the principal and the agent sign a total of $|\mathcal{A}|$ contracts, each with different recommended action $a \in \mathcal{A}$ and compensation schedule $c(a)$ as a function of the recommended action, $a$. Then, the principal and the agent would go to an authority or a third party to conduct a lottery with probability distribution function $\pi(a)$ on which contract would be implemented on that day. If the $i$-th contract is drawn from the lottery, then the third party would inform both the principal and the agent that the recommended action for that day is $a_{i}$ with the compensation schedule $c\left(a_{i}\right)$.

Arnott and Stiglitz [1] use ex ante randomization for action lotteries. This terminology refers to the situation that random contract occurs before the recommended action is chosen. They demonstrate that the action lotteries will result in a welfare improvement if the principal's expected utility is nonconcave in the agent's expected utility. However, it is not clear what the sufficient conditions would be needed for the statement in the assumption to be true.

## An NLP Formulation

When the principal proposes a contract with action lotteries, the contract has to satisfy the participation constraint and the incentive compatibility constraints. In particular, for a given contract $(\pi(a), c(a))_{a \in \mathcal{A}}$, the participation constraint requires the agent's expected utility to be at least $U^{*}$ :

$$
\begin{equation*}
\sum_{a \in \mathcal{A}} \pi(a) U(c(a), a) \geq U^{*} \tag{12}
\end{equation*}
$$

For any recommended action $a$ with $\pi(a)>0$, it has to be incentive compatible with respect to the corresponding compensation schedule $c(a) \in R^{N}$. Hence, the incentive compatibility constraints are:

$$
\begin{equation*}
\forall a \in\{\hat{a}: \pi(\hat{a})>0\}: \quad a=\operatorname{argmax}\{U(c(a), \tilde{a}): \tilde{a} \in \mathcal{A}\}, \tag{13}
\end{equation*}
$$

or equivalently,

$$
\text { if } \pi(a)>0, \text { then }\left\{\begin{array}{c}
U(c(a), a) \geq U\left(c(a), a_{1}\right)  \tag{14}\\
\vdots \\
U(c(a), a) \geq U\left(c(a), a_{M}\right)
\end{array}\right.
$$

However, we do not know in advance that at an optimal solution, whether $\pi(a)$ will be strictly positive for an action $a$. One way to overcome this difficulty is to reformulate the solution-dependent constraints (14) by:

$$
\forall a \in \mathcal{A}:\left\{\begin{array}{c}
\pi(a) U(c(a), a) \geq \pi(a) U\left(c(a), a_{1}\right)  \tag{15}\\
\vdots \\
\pi(a) U(c(a), a) \geq \pi(a) U\left(c(a), a_{M}\right)
\end{array}\right.
$$

or in a compact presentation,

$$
\begin{equation*}
\pi(a)(U(c(a), a)-U(c(a), \tilde{a})) \geq 0, \quad \forall(a, \tilde{a}(\neq a)) \in \mathcal{A} \times \mathcal{A} \tag{16}
\end{equation*}
$$

Finally, since $\pi(\cdot)$ is a probability distribution function, we need

$$
\begin{align*}
& \sum_{a \in \mathcal{A}} \pi(a)=1  \tag{17}\\
& \pi(a) \geq 0, \quad \forall a \in \mathcal{A}
\end{align*}
$$

The principal chooses a contract with action lotteries that satisfies participation constraint (12), incentive compatibility constraints (16), and the probability measure constraint (17) to maximize his expected utility. An optimal contract with action lotteries $\left(\pi^{*}(a), c^{*}(a)\right)_{a \in \mathcal{A}}$ is then a solution to the following nonlinear program:

$$
\begin{align*}
\operatorname{maximize} & \sum_{a \in \mathcal{A}} \pi(a) W(c(a), a) \\
\text { subject to } & \sum_{a \in \mathcal{A}} \pi(a) U(c(a), a) \geq U^{*}, \\
& \sum_{a \in \mathcal{A}} \pi(a)=1,  \tag{18}\\
& \pi(a)(U(c(a), a)-U(c(a), \tilde{a})) \geq 0, \quad \forall(a, \tilde{a}(\neq a)) \in \mathcal{A} \times \mathcal{A}, \\
& \pi(a) \geq 0, \quad \forall a \in \mathcal{A} .
\end{align*}
$$

The nonlinear program (18) has $(|\mathcal{Q}| *|\mathcal{A}|+|\mathcal{A}|)$ variables and $(|\mathcal{A}| *(|\mathcal{A}|-1)+2)$ constraints. In addition, its feasible region is highly nonconvex due to the last two sets of constraints in (14). As shown in the following graph, the feasible region $\mathcal{F}_{2}=\{(x, y) \mid x y \geq 0, x \geq 0\}$ is the union of the first quadrant and the $y$-axis. Furthermore, the standard nonlinear programming constraint qualification fails to hold at every point on the $y$-axis.


Fig. 2. The feasible region $\mathcal{F}_{2}=\{(x, y) \mid x y \geq 0, x \geq 0\}$.

## An MPEC formulation

Another formulation of the incentive compatibility constraints (14) is to introduce a binary variable $y(a) \in\{0,1\}$ for each action $a \in \mathcal{A}$. Then the condition (14) and nonnegativity
constraint on $\pi(a)$ can be replaced by the following complementarity constraints:

$$
\forall a \in \mathcal{A}:\left\{\begin{array}{l}
(1-y(a))(U(c(a), a)-U(c(a), \tilde{a})) \geq 0, \quad \forall \tilde{a}(\neq a) \in \mathcal{A}  \tag{19}\\
0 \leq \pi(a) \perp y(a) \geq 0 \\
y(a) \in\{0,1\}
\end{array}\right.
$$

After replacement of the last two sets of constraints in (18) by (19), the resulting MPEC, with variables $(\pi(a), c(a), y(a))_{a \in \mathcal{A}}$, for the optimal contract with action lotteries problem is:

$$
\begin{array}{ll}
\operatorname{maximize} & \sum_{a \in \mathcal{A}} \pi(a) W(c(a), a) \\
\text { subject to } & \sum_{a \in \mathcal{A}} \pi(a) U(c(a), a) \geq U^{*}, \\
& \sum_{a \in \mathcal{A}} \pi(a)=1,  \tag{20}\\
\forall a \in \mathcal{A}: \quad\left\{\begin{array}{l}
(1-y(a))(U(c(a), a)-U(c(a), \tilde{a})) \geq 0, \quad \forall \tilde{a}(\neq a) \in \mathcal{A}, \\
0 \leq \pi(a) \perp y(a) \geq 0, \\
y(a) \in\{0,1\} .
\end{array}\right.
\end{array}
$$

Allowing the compensation schedules to be dependent on the agent's action will increase the principal's expected utility; see Theorem 2. The difference between the optimal objective value of the NLP (18) (or the MPEC(20)) and that of the MPEC (11) characterizes the principal's improved welfare from using an optimal contract action lotteries.

Theorem 2 The principal prefers an optimal contract with action lotteries to an optimal contract. His expected utility from choosing an optimal contract with action lotteries will be at least as good as that of choosing an optimal contract.

### 3.2 The Contract with Compensation Lotteries

Definition 4 For any outcome $q \in \mathcal{Q}$, a randomized compensation $\tilde{c}(q)$ is a random variable on the set of compensations $\mathcal{C}$ with a probability measure $F(\cdot)$.

Remark If the set of compensations $\mathcal{C}$ is a closed interval $[\underline{c}, \bar{c}] \in R$, then the measure of $\tilde{c}(q)$ is a cumulative density function (cdf) $F:[\underline{c}, \bar{c}] \rightarrow[0,1]$ with $F(\underline{c})=0$ and $F(\bar{c})=1$. In addition, $F(\cdot)$ is nondecreasing and right-continuous.

To simplify the analysis, we assume that every randomized compensation $\tilde{c}(q)$ has finite support.

Assumption 1: Finite support for randomized compensation. For all $q \in \mathcal{Q}$, the randomized compensation $\tilde{c}(q)$ has finite support over an unknown set $\left\{c_{1}(q), c_{2}(q), \ldots, c_{L}(q)\right\}$ with a known $L$.

An immediate consequence of Assumption 1 is that we can write $\tilde{c}(q)=c_{i}(q)$ with probability $p_{i}(q)>0$ for all $i=1, \ldots, L$ and $q \in \mathcal{Q}$. In addition, we have $\sum_{i=1}^{L} p_{i}(q)=1$ for all $q \in \mathcal{Q}$. Notice that both $\left(c_{i}(q)\right)_{i=1}^{L} \in R^{L}$ and $\left(p_{i}(q)\right)_{i=1}^{L} \in R^{L}$ are endogenous variables and will be chosen by the principal.

Definition 5 A compensation lottery is a randomized compensation schedule $\tilde{c}=\left(\tilde{c}\left(q_{1}\right), \ldots, \tilde{c}\left(q_{N}\right)\right) \in R^{N}$, in which $\tilde{c}(q)$ is a randomized compensation satisfying Assumption 1 for all $q \in \mathcal{Q}$.

Definition 6 A contract with compensation lotteries consists of a recommended action $a$ to the agent and a randomized compensation schedule $\tilde{c}=\left(\tilde{c}\left(q_{1}\right), \ldots, \tilde{c}\left(q_{N}\right)\right) \in R^{N}$.

Let $c^{q}=\left(c_{i}(q)\right)_{i=1}^{L} \in R^{L}$ and $p^{q}=\left(p_{i}(q)\right)_{i=1}^{L} \in R^{L}$. Given the outcome $q$ is observed by the principal, we let $\mathbf{w}\left(c^{q}, p^{q}\right)$ denote the principal's expected utility with respect to a randomized compensation $\tilde{c}(q)$, i.e.,

$$
\mathbf{w}\left(c^{q}, p^{q}\right)=\mathbb{E} w(q-\tilde{c}(q))=\sum_{i=1}^{L} p_{i}(q) w\left(q-c_{i}(q)\right) .
$$

With a randomized compensation schedule $\tilde{c}$ and a recommended action $a$, the principal's expected utility then becomes

$$
\begin{equation*}
\mathbb{E} W(\tilde{c}, a)=\sum_{q \in \mathcal{Q}} p(q \mid a)\left(\sum_{i=1}^{L} p_{i}(q) w\left(q-c_{i}(q)\right)\right)=\sum_{q \in \mathcal{Q}} p(q \mid a) \mathbf{w}\left(c^{q}, p^{q}\right) \tag{21}
\end{equation*}
$$

Similarly, given a recommended action $a$, we let $\mathbf{u}\left(c^{q}, p^{q}, a\right)$ denote the agent's expected utility with respect to $\tilde{c}(q)$ for the observed outcome $q$ :

$$
\mathbf{u}\left(c^{q}, p^{q}, a\right)=\mathbb{E} u(\tilde{c}(q), a)=\sum_{i=1}^{L} p_{i}(q) u\left(c_{i}(q), a\right)
$$

The agent's expected utility with a randomized compensation schedule $\tilde{c}$ and a recommended action $a$ is

$$
\begin{equation*}
\mathbb{E} U(\tilde{c}, a)=\sum_{q \in \mathcal{Q}} p(q \mid a)\left(\sum_{i=1}^{L} p_{i}(q) u\left(c_{i}(q), a\right)\right)=\sum_{q \in \mathcal{Q}} p(q \mid a) \mathbf{u}\left(c^{q}, p^{q}, a\right) \tag{22}
\end{equation*}
$$

To further simply to notation, we use $c_{\mathcal{Q}}=\left(c^{q}\right)_{q \in \mathcal{Q}}$ and $p_{\mathcal{Q}}=\left(p^{q}\right)_{q \in \mathcal{Q}}$ to denote the collection of variables $c^{q}$ and $p^{q}$, respectively. We also let $\mathbf{W}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a\right)$ denote the principal's expected utility $\mathbb{E} W(\tilde{c}, a)$ as defined in $(21)$, and similarly, $\mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a\right)$ for $\mathbb{E} U(\tilde{c}, a)$ as in (22).

An optimal contract with compensation lotteries $\left(c_{\mathcal{Q}}^{*}, p_{\mathcal{Q}}^{*}, a^{*}\right)$ is a solution to the following problem:

$$
\begin{array}{cl}
\operatorname{maximize} & \mathbf{W}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a\right) \\
\text { subject to } & \mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a\right) \geq U^{*} \\
& \mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a\right) \geq \mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a_{1}\right)  \tag{23}\\
& \vdots \\
& \mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a\right) \geq \mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a_{M}\right), \\
& a \in \mathcal{A}=\left\{a_{1}, \ldots, a_{M}\right\} .
\end{array}
$$

Define

$$
\begin{aligned}
\mathbf{W}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}\right) & =\left(\mathbf{W}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a_{1}\right), \ldots, \mathbf{W}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a_{M}\right)\right) \in R^{M} \\
\mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}\right) & =\left(\mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a_{1}\right), \ldots, \mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a_{M}\right)\right) \in R^{M}
\end{aligned}
$$

Following the derivation as in Section 2, we can reformulate the program for an optimal contract with compensation lotteries (23) as a mixed-integer nonlinear program with decision variables $\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}\right)$ and $y=\left(y_{i}\right)_{i=1}^{M}$ :

$$
\begin{array}{cl}
\text { maximize } & y^{\mathrm{T}} \mathbf{W}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}\right) \\
\text { subject to } & y^{\mathrm{T}} \mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}\right) \geq U^{*}, \\
& y^{\mathrm{T}} \mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}\right) \geq \mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a_{1}\right), \\
& \vdots  \tag{24}\\
& y^{\mathrm{T}} \mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}\right) \geq \mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}, a_{M}\right), \\
& e_{M}^{\mathrm{T}} y=1, \\
& y_{i} \in\{0,1\} \quad \forall i=1, \ldots, M,
\end{array}
$$

Similarly, the MPEC formulation with decision variables $\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}\right)$ and $\delta \in R^{M}$ is:

$$
\begin{array}{ll}
\operatorname{maximize} & \delta^{\mathrm{T}} \mathbf{W}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}\right) \\
\text { subject to } & \delta^{\mathrm{T}} \mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}\right) \geq U^{*} \\
& e_{M}^{\mathrm{T}} \delta=1  \tag{25}\\
& 0 \leq \delta \perp\left(\delta^{\mathrm{T}} \mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}\right)\right) e_{M}-\mathbf{U}\left(c_{\mathcal{Q}}, p_{\mathcal{Q}}\right) \geq 0
\end{array}
$$

Arnott and Stiglitz [1] call the compensation lotteries ex post randomization; this refers to the situation that the random compensation occurs after the recommended action is chosen or implemented. They show that if the agent is risk averse and his utility function is separable, and if the principal is risk neutral, then the compensation lotteries are not desirable.

### 3.3 The Contract with Action and Compensation Lotteries

Definition 7 A contract with action and compensation lotteries is a probability distribution over actions, $\pi(a)$, and a randomized compensation schedule $\tilde{c}(a)=\left(\tilde{c}\left(q_{1}, a\right), \ldots, \tilde{c}\left(q_{N}, a\right)\right) \in$ $R^{N}$ for every $a \in \mathcal{A}$; The randomized compensation schedule $c(a)$ is an agreement between the principal and the agent such that $\tilde{c}(q, a) \in \mathcal{C}$ is a randomized compensation to the agent from the principal if outcome $q \in \mathcal{Q}$ is observed and the action $a \in \mathcal{A}$ is recommended by the principal.

Assumption 2 For every action $a \in \mathcal{A}$, the randomized compensation schedule $\tilde{c}(q, a)$ satisfies the finite support assumption (Assumption 1) for all $q \in \mathcal{Q}$.

With Assumption 2, the notation $c^{q}(a), p^{q}(a), c_{\mathcal{Q}}(a), p_{\mathcal{Q}}(a)$ is analogous to what we have defined in Section 3.1 and 3.2. Without repeating the same derivation process described earlier, we give the MPEC formulation for the optimal contract with action and compensation lotteries problem with variables $\left(\pi(a), c_{\mathcal{Q}}(a), p_{\mathcal{Q}}(a), y(a)\right)_{a \in \mathcal{A}}$ :

$$
\begin{array}{ll}
\operatorname{maximize} & \sum_{a \in \mathcal{A}} \pi(a) \mathbf{W}\left(c_{\mathcal{Q}}(a), p_{\mathcal{Q}}(a), a\right) \\
\text { subject to } & \sum_{a \in \mathcal{A}} \pi(a) \mathbf{U}\left(c_{\mathcal{Q}}(a), p_{\mathcal{Q}}(a), a\right) \geq U^{*}, \\
& \sum_{a \in \mathcal{A}} \pi(a)=1 \\
\forall a \in \mathcal{A}: \quad\left\{\begin{array}{l}
\forall \tilde{a}(\neq a) \in \mathcal{A}: \\
(1-y(a))\left(\mathbf{U}\left(c_{\mathcal{Q}}(a), p_{\mathcal{Q}}(a), a\right)-\mathbf{U}\left(c_{\mathcal{Q}}(a), p_{\mathcal{Q}}(a), \tilde{a}\right)\right) \geq 0, \\
0 \leq \pi(a) \perp y(a) \geq 0, \\
y(a) \in\{0,1\} .
\end{array}\right. \tag{26}
\end{array}
$$

## Linear Programming Approximation

Townsend [24, 25] was among the first to use linear programming techniques to compute static incentive constrained problems. Prescott [18, 19] further apply linear programming specifically to solve moral-hazard problems. A solution obtained by the linear programming approach is an approximation to a solution to the $\operatorname{MPEC}(26)$. Instead of treating $c_{\mathcal{Q}}(a)$ as unknown variables, one can construct a grid $\Xi$ with element $\xi$ to approximate the set $\mathcal{C}$ of compensations. By introducing probability measures associated with the action lotteries on $\mathcal{A}$ and compensation lotteries on $\Xi$, one can then approximate a solution to the moral-hazard problem with lotteries (26) by solving a linear program. More specifically, the principal chooses probability distributions $\pi(a)$, and $\pi(\xi \mid q, a)$ over the set of actions $\mathcal{A}$, the set of outcomes $\mathcal{Q}$, and the compensation grid $\Xi$. One then can reformulate the resulting nonlinear
program into a linear program with decision variables $\pi=(\pi(\xi, q, a))_{\xi \in \Xi, q \in \mathcal{Q}, a \in \mathcal{A}}$ :

$$
\begin{align*}
\operatorname{maximize}_{(\pi)} & \sum_{\xi, q, a} w(q-\xi) \pi(\xi, q, a) \\
\text { subject to } & \sum_{\xi, q, a} u(\xi, a) \pi(\xi, q, a) \geq U^{*}, \\
& \sum_{\xi, q} u(\xi, a) \pi(\xi, q, a) \geq \sum_{\xi, q} u(\xi, \tilde{a}) \frac{p(q \mid \tilde{a})}{p(q \mid a)} \pi(\xi, q, a) \forall(a, \tilde{a}(\neq a)) \in \mathcal{A} \times \mathcal{A},  \tag{27}\\
& \sum_{\xi} \pi(\xi, \tilde{q}, \tilde{a})=p(\tilde{q} \mid \tilde{a}) \sum_{\xi, q} \pi(\xi, q, \tilde{a}) \quad \forall(\tilde{q}, \tilde{a}) \in \mathcal{Q} \times \mathcal{A}, \\
& \sum_{\xi, q, a} \pi(\xi, q, a)=1, \\
& \pi(\xi, q, a) \geq 0 \quad \forall(\xi, q, a) \in \Xi \times \mathcal{Q} \times \mathcal{A} .
\end{align*}
$$

Note that the above linear program has $(|\Xi| *|\mathcal{Q}| *|\mathcal{A}|)$ variables and $(|\mathcal{A}| *(|\mathcal{Q}|+|\mathcal{A}|-1)+2)$ constraints. The size of the linear program will grow enormously and quickly when one chooses a fine grid. For example, if there are 50 actions, 40 outputs, and 500 compensations, then the linear program has one million variables and 4452 constraints. It will become computationally infeasible because of the limitation on the computer memory, if not the time required to solve a large-scale linear program. On the other hand, a solution of the LP obtained from a coarse grid will not be satisfactory if an accurate solution is needed. Prescott [19] points out that the constraint matrix of the linear program (27) has block angular structure. As a consequence, one can apply the Dantzig-Wolfe decomposition algorithm to the linear program (27) to lessen the need for computer memory and the computational time. Recall that the MPEC (11) for the optimal contract problem has only $(|\mathcal{Q}|+|\mathcal{A}|)$ variables and $|\mathcal{A}|$ complementarity constraints with one linear constraint and one nonlinear constraint. Even with the use of the Dantzig-Wolfe decomposition algorithm to solve the LP (27), choosing the "right" grid is still an issue. With the advances in both theory and numerical methods for solving MPECs in the last decade, we believe that the MPEC approach has greater advantages in solving a much smaller problem and in obtaining a more accurate solution.

The error from discretizing set of compensations $\mathcal{C}$ is characterized by the difference between the optimal objective value of the LP (27) and that of the MPEC (26).

Theorem 3 Assume the agent is risk averse over the payoff received from the principal and his effort and the principal is risk averse or risk neutral over his net income. The feasible region of the LP (27) is contained in the feasible region of the MPEC (26). The optimal objective value of the LP (27) is lower than that of the MPEC (26).

## 4 A Hybrid Approach toward Global Optimization

One reason that nonconvex programs are not popular among economists is the issue of the need for global solutions. While the local search algorithms for solving nonconvex programs have fast convergence properties near a solution, they are designed to find a local solution. Algorithms for solving MPECs are no exception. One heuristic in practice is to solve the same problem with several different starting points. It then becomes a trade off between the computational time and the quality of the "best" solution found.

Linear programming does not suffer from the global solution issue. However, to obtain an accurate solution to a moral-hazard problem via the linear programming approach, one needs to use a very fine compensation grid. This often leads to solve large-scale linear programs with millions of variables and ten or hundred thousands of constraints, which might be computationally infeasible, due to insufficient computer memory and long computational time.

Certainly, there is a need to develop a global optimization method with fast local convergence for MPECs. Below, we propose a hybrid approach combining both MPECs and linear programming approaches to find a global solution (or at least better than the LP solution) of an optimal contract problem. The motivation for this hybrid method comes from the observation that the optimal objective value of the LP approach from a coarse grid could provide a lower bound on the optimal objective value of the MPEC as well as a good guess on the final recommended action $a^{*}$. We then can use this information to exclude some undesired local minimizers and to provide a good starting point when we solve the MPEC (11). This heuristic procedure toward a global solution of the MPEC (11) leads to the following algorithm.

## A hybrid method for the optimal contract problem as an MPEC (11)

Step 0: Construct a coarse grid.
Step 1: Solve the LP (27) for the given grid.
Step 2: $\begin{cases}(2.1): & \text { Compute } \pi(\cdot, \cdot, a)=\sum_{\xi \in \Xi} \sum_{q \in \mathcal{Q}} \pi(\xi, q, a), \quad \forall a \in \mathcal{A} ; \\ (2.2): & \text { Compute } \mathbb{E}[\xi(q)]=\sum_{\xi \in \Xi} \xi \pi(\xi, q, a), \quad \forall q \in \mathcal{Q} ; \\ (2.3): & \text { Set the initial point } c^{0}=(\mathbb{E}[\xi(q)])_{q \in \mathcal{Q}} \text { and } \delta^{0}=(\pi(\cdot, \cdot, a))_{a \in \mathcal{A}} ; \\ (2.4): & \text { Solve the MPEC (11) with the starting point }\left(c^{0}, \delta^{0}\right) .\end{cases}$
Step 3: Refine the grid and repeat Step 1 and Step 2.

Remark If the starting point from an LP solution is close to the optimal solution of the MPEC (11), then the sequence of iterates generated by the SQP algorithm converge Q-
quadratically to the optimal solution. See Proposition 5.2 in Fletcher et al. [5].
One can also develop the similar procedures to solve the global solutions for optimal contract problems with action and/or compensation lotteries. However, the MPECs for contracts with lotteries are much more numerically challenging problems than the MPEC (11) for deterministic contracts.

## 5 An Example and Numerical Results

To illustrate the use of the mixed-integer nonlinear program (8), the MPEC (11) and the hybrid approaches, and to understand the effect on discretizing the set of compensations $\mathcal{C}$, we only consider to problems of deterministic contracts without lotteries. We consider a two-outcome example in Karaivanov [10]. Before starting the computational work, we summarize the problem characteristics of various approaches to compute the optimal deterministic contracts in Table 5.1.

Table 5.1. Problem characteristics of various approaches

|  | MINLP (8) | MPEC (11) | LP $(27)$ |
| :--- | :---: | :---: | :---: |
| $\#$ of Variables | $\|\mathcal{Q}\|$ | $\|\mathcal{Q}\|+\|\mathcal{A}\|$ | $\|\Xi\| *\|\mathcal{Q}\| *\|\mathcal{A}\|$ |
| $\#$ of Binary Variables | $\|\mathcal{A}\|$ | - | - |
| $\#$ of Constraints | $\|\mathcal{A}\|+2$ | 2 | $\|\mathcal{A}\| *(\|\mathcal{Q}\|+\|\mathcal{A}\|-1)+2$ |
| $\#$ of Complem. Constraints | - | $\|\mathcal{A}\|$ | - |

## Example: No Action and Compensation Lotteries

Assume the principal is risk neutral with utility $w(q-c(q))=q-c(q)$, and the agent is risk averse with utility $u(c(q), a)=\frac{c^{1-\gamma}}{1-\gamma}+\kappa \frac{(1-a)^{1-\delta}}{1-\delta}$. Suppose there are only two possible outcomes, e.g., a coin-flip. If the desirable outcome (high sale quantities or high production quantities) happens, then the principal receives $q_{H}=\$ 3$; otherwise, he receives $q_{L}=\$ 1$. For simplicity, we assume that set of actions $\mathcal{A}$ consists of $|\mathcal{A}|$ equally-spaced effort level within the closed interval $[0.01,0.99]$. The production technology for the high outcome is described by $p\left(q=q_{H} \mid a\right)=a^{\alpha}$ with $0<\alpha<1$. Notice that since 0 and 1 are excluded from the action set $\mathcal{A}$, the full-support assumption on production technology is satisfied.

The value of the parameters for the particular instance we solve is given in Table 5.2.

Table 5.2. The value of parameters used in the example.

| $\gamma$ | $\kappa$ | $\delta$ | $\alpha$ | $U^{*}$ | $\|\mathcal{A}\|$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0.5 | 1 | 0.5 | 0.7 | 1 | 10 |

We solve this problem first as an mixed-integer nonlinear program (8) and then as an MPEC (11). For the LP lotteries approach, we start with 20 grid points in the compensation grid (we evenly discretize the compensation set $\mathcal{C}$ into 19 segments) and then increase the size of the compensation grid to $50,100,200, \cdots, 5000$.

We submitted the corresponding AMPL programs to the NEOS server [15]. The mixedinteger nonlinear programs were solved using the MINLP solver [3] on the computer host newton.mcs.anl.gov. To obtain fair comparisons between the LP, MPEC, and hybrid approaches, we chose SNOPT [6] to solve the associated mathematical programs. The AMPL programs were solved on the computer host tate.iems.northwestern.edu.

Table 5.3 gives the solutions returned by the MINLP solver to the mixed-integer nonlinear program (8). We use $y=0$ and $y=e_{M}$ as starting points. In both cases, the MINLP solver returns a solution very fast. However, it does not guarantee to find a global solution.

Table 5.3. Solutions of the MINLP approach.

| Starting <br> Point | \# of Regular <br> Variables | \# of Binary <br> Variables | \# of . <br> Constr. | Solve Time <br> (in sec.) | Objective <br> Value |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $y=0$ | 2 | 10 | 12 | 0.01 | 1.864854251 |
| $y=e_{M}$ | 2 | 10 | 12 | 0.00 | 1.877265189 |

For solving the MPEC (11), we try two different starting points to illustrate the possibility of finding only a local solution. The MPEC solutions are given in Table 5.4 below.

Table 5.4. Solutions of the MPEC approach with two different starting points.

| Starting <br> Point | \# of <br> Variables | \# of Complem. <br> Constraints | Read Time <br> (in sec.) | Solve Time <br> (in sec.) | \# of Major <br> Iterations | Objective <br> Value |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\delta=0$ | 22 | 10 | 0 | 0.07 | 45 | 1.079621424 |
| $\delta=e_{M}$ | 22 | 10 | 0 | 0.18 | 126 | 1.421561553 |

The solutions for the LP lottery approach with different compensation grids are given in Table 5.5. Notice that the solve time increases faster than the size of the grid when $|\Xi|$ is in the order of $10^{5}$ and higher, while the number of major iterations only increase about 3 times when we increase the grid size 250 times (from $|\Xi|=20$ to $|\Xi|=5000$ ).

Table 5.5. Solutions of the LP approach with 8 different grids (\# of constraints $=112$ ).

| $\|\Xi\|$ | \# of <br> Variables | Read Time <br> (in sec.) | Solve Time <br> (in sec.) | \# of Major <br> Iterations | Objective <br> Value |
| ---: | :---: | :---: | :---: | :---: | :---: |
| 20 | 400 | 0.01 | 0.03 | 31 | 1.876085819 |
| 50 | 1000 | 0.02 | 0.06 | 46 | 1.877252488 |
| 100 | 2000 | 0.04 | 0.15 | 53 | 1.877252488 |
| 200 | 4000 | 0.08 | 0.31 | 62 | 1.877254211 |
| 500 | 10000 | 0.21 | 0.73 | 68 | 1.877263962 |
| 1000 | 20000 | 0.40 | 2.14 | 81 | 1.877262184 |
| 2000 | 40000 | 0.83 | 3.53 | 71 | 1.877260460 |
| 5000 | 100000 | 2.19 | 11.87 | 101 | 1.877262793 |

Finally, for the hybrid approach, we first use the LP solution from a compensation grid with $|\Xi|=20$ to construct a starting point for the MPEC (11). As one can see in Table 5.6 , with a good starting point, it only takes SNOPT 0.01 second to find a solution to the example formulated as the MPEC (11). Furthermore, the optimal objective value is higher that of the LP solution from a fine compensation grid with $|\Xi|=5000$.

Table 5.6. Solution of the hybrid approach.

| $\mathbf{L P}$ <br> $\|\Xi\|$ | Read Time <br> (in sec.) | Solve Time <br> (in sec.) | \# of Major <br> Iterations | Objective <br> Value |
| :---: | :---: | :---: | :---: | :---: |
| 20 | 0.01 | 0.03 | 31 | 1.876085819 |
| MPEC <br> Starting Point | Read Time <br> (in sec.) | Solve Time <br> (in sec.) | \# of Major <br> Iterations | Objective <br> Value |
| $\delta_{6}=1, \delta_{i(\neq 6)}=0$ | 0.02 | 0.01 | 13 | 1.877265298 |

## 6 Conclusions and Future Work

The purpose of this paper is to introduce the MPEC approach and apply it to moralhazard problems. We have presented MPEC formulations for optimal deterministic contract problems and optimal contract problems with action and/or compensation lotteries. We also formulated the former problem as a mixed-integer nonlinear program. To obtain a global solution, we have proposed a hybrid procedure that combines the LP lottery and the MPEC approaches. In this procedure, the LP solution from a coarse compensation grid provides a good starting point for the MPEC. We then can apply specialized MPEC algorithms with fast local convergence rate to obtain a solution. In an numerical example, we have demonstrated that the hybrid method is more efficient than using only the LP lottery approach, which requires to solve a sequence of large-scale linear programs. Although we can not prove that the hybrid approach will guarantee to find a global solution, it finds one better than the solution from the LP lottery approach. We plan to test the numerical performance of the
hybrid procedure on other examples such as the bank regulation example in [18] and the two-dimensional action choice example in [19].

One can extend the MPEC approach to single-principal multiple-agent problems without any problem. For the model of multiple-principal multiple-agent [14], it can be formulated as an equilibrium problem with equilibrium constraint (EPEC) [23]. We will investigate these two topics in our future research.

Another important topic we plan to explore is the dynamic moral-hazard problem; see Phelan and Townsend [17]. In the literature, dynamic programming is applied to solve this model. We believe that there is an equivalent nonlinear programming formulation. Analogous to the hybrid procedure proposed in Section 4, an efficient method to solve this dynamic model is to combine dynamic programming and nonlinear programming.

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